

# Research Report 396

# THE IMPACT OF PREDICTIVE INACCURACIES ON EXECUTION SCHEDULING

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# The Impact of Predictive Inaccuracies on Execution Scheduling\*

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Abstract: This paper investigates the underlying impact of predictive inaccuracies on

execution scheduling, with particular reference to execution time predictions. This

study is conducted from two perspectives: from that of job selection and from that of

resource allocation, both of which are fundamental components in execution sched-

uling. A new performance metric, termed the degree of misperception, is introduced

to express the probability that the predicted execution times of jobs display different

ordering characteristics from their real execution times due to inaccurate prediction.

Specific formulae are developed to calculate the degree of misperception in both job

selection and resource allocation scenarios. The parameters which influence the de-

gree of misperception are also extensively investigated. The results presented in this

paper are of significant benefit to scheduling approaches that take into account pre-

dictive data; the results are also of importance to the application of these scheduling

techniques to real-world high-performance systems.

Index terms: performance prediction, execution time, scheduling, job selection, re-

source allocation, and performance evaluation

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#### 1. Introduction

Scheduling in a single processor environment is, at its most basic level, the task of determining the sequence in which jobs should be executed. In a multi-processor or multi-computer environment on the other hand, job scheduling also involves the process of resource allocation, that is, determining the resource to which a job should be sent for execution. The design of scheduling policies for parallel and distributed systems has received a good deal of attention [2, 4, 5, 14, 15]. These schemes are often based on the assumption that the job execution times are known somehow [5, 9]. While this is a useful assumption to make, information of this type is often unknown and must therefore be obtained through some predictive mechanism. A naïve approach to this problem might be to require the owner of the job to estimate the execution time based on their past experience. A more sophisticated approach might be to utilize performance prediction tools to perform this function. A number of increasingly accurate prediction tools have been developed that are able to predict the execution times of jobs based on performance models [3, 6, 7, 8] or historical data [1, 13].

In spite of this, an inevitable fact is that the prediction results are unlikely to be entirely accurate and as a result, this may have a fundamental impact on job selection and resource allocation.

In the case of job selection, the inaccurate prediction may mean that the scheduler has an incorrect view of the order in which the different jobs should execute. For example, it may be the case that the real execution time of job  $J_I$  is greater than that of job  $J_2$ ; because of the inaccurate prediction however, the scheduler may view job  $J_I$  as having a shorter execution time than that of job  $J_2$ . Therefore, if the scheduling policy is based on job execution times (the shortest job serviced first, for example), then this

misperception will impact on the order in which jobs are selected for execution. This will ultimately influence the scheduler and system performance.

When a scheduler receives a job in a parallel or distributed system, there may be a number of resources (processors or computers) available on which the job may be executed. If the resource allocation policy is also based on the expected execution time of the job on the different resources (select the computer that offers the shortest execution time, for example) then these inaccuracies might also cause the scheduler to wrongly select between them. Again, this misperception will impact on the scheduling and overall system performance.

The existence of misperception originates from the inaccurate prediction (of in this case execution time), and should be viewed as an inherent characteristic of any prediction-based scheduling scheme that operates in a highly-variable real-world system. This said, different scheduling policies will have different levels of sensitivity to the degree of misperception. Thus, the study of the impact of inaccurate prediction on scheduling performance can be decomposed into two levels – at the underlying level, as a study into the degree of misperception originating from inaccurate prediction; and at the higher level, as a study into the sensitivity of individual scheduling policies to this degree of misperception. This paper addresses the former, where the latter is the subject of future work.

Different predicted errors can lead to different degrees of misperception. The objective of this paper is to establish the relationship between the predicted error and the degree of misperception, for both job selection and resource allocation scenarios. This study provides an insight into the underlying impact of inaccurate prediction on job selection and resource allocation and significantly benefits the design and evaluation of scheduling policies that use predictive data.

The rest of this paper is organized as follows. Formula for the degree of misperception, and for job selection and resource allocation are developed in Section 2. The parameters that influence the degree of misperception are then extensively evaluated in Section 3. The paper concludes in Section 4.

## 2. An Analysis of the Degree of Misperception

#### 2.1 Job Selection

When performance prediction tools are used to estimate the execution times of jobs, the predicted execution time usually lies in an interval around the actual execution time (of the job) according to some probability distribution [7, 8].

Suppose that the actual execution time of job  $J_i$  is  $x_i$  and that the predictive error, denoted by  $y_i$ , is a random variable in the range  $[-ax_i, bx_i]$  following some probability density function,  $g_i(y_i)$ ; where the possible value fields of a and b are [0, 100%] and  $[0, \infty)$ , respectively. It is assumed that the predictive errors of different jobs are independent random variables. The predicted execution time of job  $J_i$ , denoted by  $z_i$ , is computed using Eq.1.

$$z_i = x_i + y_i \tag{1}$$

Suppose two jobs  $J_I$  and  $J_2$  have the actual execution times  $x_I$  and  $x_2$ , where  $x_I < x_2$ . Then the predicted execution times of  $J_I$  and  $J_2$ , that is  $z_I$  and  $z_2$ , are  $x_I + y_I$  and  $x_2 + y_2$ , respectively. Given that  $x_I < x_2$ , a misperception happens if  $z_I \ge z_2$ . The degree of misperception for these two jobs, denoted by  $MD(x_I, x_2)$ , is defined by the probability that  $z_I \ge z_2$  while  $x_I < x_2$ . This probability is denoted by  $P_r(z_I \ge z_2 | x_I < x_2)$ . With Eq.1, this probability can be further transformed using Eq.2.

$$P_r(z_1 \ge z_2 | x_1 < x_2) = P_r(x_1 + y_1 \ge x_2 + y_2 | x_1 < x_2) = P_r(y_1 \ge y_2 + x_2 - x_1 | x_2 - x_1 > 0)$$
 so that the  $MD(x_1, x_2)$  is computed by Eq.3. (2)

$$MD(x_1, x_2) = P_r(y_1 \ge y_2 + x_2 - x_1 | x_2 - x_1 > 0)$$
(3)

From Eq.3, it can be seen that the probability that the misperception happens is the probability that given  $x_2-x_1>0$ , the predictive error of  $J_1$  (i.e.,  $y_1$ ) is greater than the predictive error of  $J_2$  (i.e.,  $y_2$ ) plus the difference between  $x_2$  and  $x_1$ . By constructing the coordinates of the predictive error  $y_1$  and  $y_2$ , the inequality  $y_1 \ge y_2 + x_2 - x_1$  means that  $y_1$  and  $y_2$  are given values from the area above the line  $y_1 = y_2 + x_2 - x_1$ .

Fig.1.a, 2.a and 3.a illustrate the relation of predicted execution times of  $J_1$  and  $J_2$ , while Fig.1.b, 2.b and 3.b show the corresponding value fields of the predictive error of  $J_1$  and  $J_2$  ( $y_1$  and  $y_2$ ) as well as the corresponding area in which  $y_1 \ge y_2 + x_2 - x_1$ .

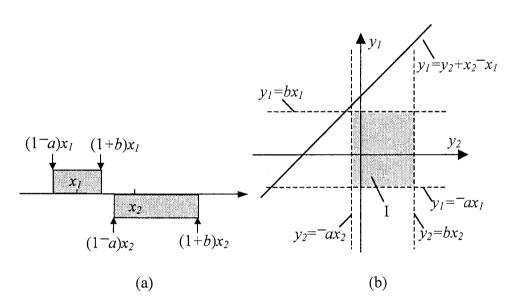


Figure.1. Case study 1 (a) the predicted execution times of jobs  $J_1$  and  $J_2$  do not overlap; (b) the corresponding coordinate area which the predictive errors  $y_1$  and  $y_2$  can be assigned values from.

Fig. 1.a illustrates the case when the ranges of predicted execution times of  $J_l$  and  $J_2$  do not overlap. In this case a misperception will not occur even if the predictions are not accurate. Correspondingly, Fig1.b shows the coordinate area of the predicted errors of  $J_l$  and  $J_2$  (area I), which is the area surrounded by the lines:  $y_l = -ax_l$ ,  $y_l = bx_l$ ,

 $y_2=-ax_2$  and  $y_2=bx_2$ . As can be seen from the figure, all of area I is below the line  $y_1=y_2+x_2-x_1$ . Hence,  $P_r(y_1 \ge y_2+x_2-x_1|x_2-x_1>0)$  in Eq.2 is equal to zero. This case is expressed more formally below.

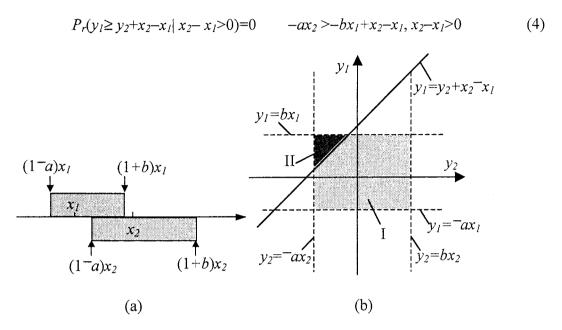


Figure.2. Case study 2 (a) the predicted execution times of jobs  $J_1$  and  $J_2$  overlap, however the lower limit of the predicted execution time for  $J_2$  does not cover  $x_1$ ; (b) the corresponding coordinate area of predicted errors of  $J_1$  and  $J_2$  ( $y_1$  and  $y_2$ ) in which the misperception occurs (the area is a triangle).

Fig. 2.a illustrates the case when the predicted execution times of  $J_l$  and  $J_2$  overlap and the lower limit of the predicted execution time of  $J_2$  is greater than  $x_l$ . The corresponding coordinate area of  $y_l$  and  $y_2$  is shown in Fig. 2.b (area I). In area I, part of the area is above the line  $y_l=y_2+x_2-x_l$  (area II), which is itself surrounded by the three lines:  $y_l=bx_l$ ,  $y_2=-ax_2$  and  $y_l=y_2+x_2-x_l$ . When  $y_l$  and  $y_2$  are assigned values from area II,  $y_l$  and  $y_2$  satisfy  $y_l \ge y_2+x_2-x_l$ , that is, a misperception will occur; a misperception will not occur if  $y_l$  and  $y_2$  are assigned values from any other area in I, although prediction errors will still exist. The probability in Eq.2 equals the double integral of the probability density functions of predicted errors (i.e.,  $g_l(y_l)$  and  $g_2(y_2)$ ) on area II, which is calculated in Eq.5.

$$P_{r}(y_{l} \ge y_{2} + x_{2} - x_{l} | x_{2} - x_{l} \ge 0) = \int_{-ax_{2} + x_{2} - x_{1}}^{bx_{1}} \int_{-ax_{2}}^{y_{1} - (x_{2} - x_{1})} g_{1}(y_{1}) g_{2}(y_{2}) dy_{2} dy_{1}$$

$$-ax_{1} - (x_{2} - x_{1}) < -ax_{2} \le bx_{1} - (x_{2} - x_{1}), x_{2} - x_{1} \ge 0$$
 (5)

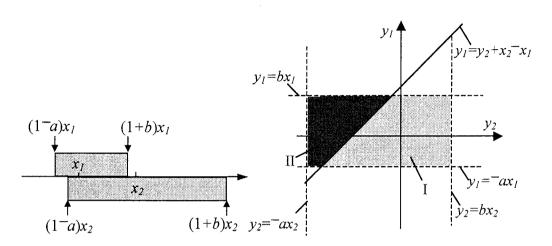


Figure.3. Case study 3 (a) the predicted execution times of jobs  $J_1$  and  $J_2$  overlap and the lower limit of predicted execution time of  $J_2$  covers  $x_1$ ; (b) the corresponding coordinate area of predicted errors of  $J_1$  and  $J_2$  ( $y_1$  and  $y_2$ ) in which the misperception occurs (the area is a trapezoid).

Fig.3.a illustrates a third case when the predicted execution times of  $J_1$  and  $J_2$  overlap. The difference from Fig.2.a is that the value field of  $J_2$ 's predicted execution time covers  $x_1$ . The corresponding coordinate area in which  $y_1$  and  $y_2$  satisfy  $y_1 \ge y_2 + x_2 - x_1$  is highlighted as area II in Fig.3.b. The area is a trapezoid rather than a triangle as in Fig.2.b. The formula for calculating the probability in Eq.2 is also different from Eq.5; this is shown in Eq.6.

$$P_{r}(y_{l} \ge y_{2} + x_{2} - x_{l} | x_{2} - x_{l} > 0) = \int_{-ax_{1}}^{bx_{1}} \int_{-ax_{2}}^{y_{1} - (x_{2} - x_{1})} g_{1}(y_{1}) g_{2}(y_{2}) dy_{2} dy_{1}$$

$$-ax_{2} \le -ax_{1} - (x_{2} - x_{1}), x_{2} - x_{1} > 0 \quad (6)$$

Eq.4-6 account for all possible relations between the predicted execution times of jobs  $J_1$  and  $J_2$ .

Suppose that the probability density function of a jobs' actual execution time in a job stream is f(x) and that the value field of the real execution time x is [xl, xu]. The degree of misperception of this job stream, denoted by  $\overline{MD}$ , is defined by the average of the degree of misperception for any two jobs in the job stream.  $\overline{MD}$  is computed using Eq.7, where  $MD(x_1, x_2)$  is computed using Eq.4-6.

$$\overline{MD} = \int_{x_1}^{x_2} \int_{x_1}^{x_2} f(x_1) f(x_2) MD(x_1, x_2) dx_2 dx_1$$
 (7)

There are several independent parameters in Eq.7, as well as in Eq.4-6: a, b, xu, xl, f(x), and  $g_i(x_i)$ . It is highly beneficial to study how these parameters influence the value of  $\overline{MD}$ ; this is the subject of the investigation in Section 3.

#### 2.2 Resource allocation

In dedicated environments, the execution time of one unit of work can be represented as a predicted point value [10, 11, 12]. However, in non-dedicated environments, the existence of background workloads on the resources causes a variation in unit execution times [3, 10, 11, 12]. Hence it can be assumed that the actual execution time of one unit of work locates across a range around the predicted point value following a certain probability [10, 12, 16].

Suppose there is a distributed system consisting of n heterogeneous computers  $c_1$ ,  $c_2$ , ...,  $c_n$ . Computer  $c_i$  is weighted  $w_i$  ( $1 \le i \le n$ ), which represents the time it takes to perform one unit of computation. Now suppose for any i, j ( $1 \le i, j \le n$ ),  $w_i < w_j$  if i < j. A job with size s is therefore predicted to have the execution time  $sw_i$  on computer  $c_i$ . The predicted execution time of a job is denoted by  $z_{ci}$ , that is,

$$z_{ci} = sw_i \tag{8}$$

However, in shared environments the actual execution time for a job with size s on computer  $c_i$  may not be  $sw_i$ , because of the existence of background workload. The actual execution time of a job on  $c_i$  is therefore denoted by  $x_{ci}$ .

For a job with size s, its predicted error on computer  $c_i$ , denoted by  $y_{ci}$ , is computed using Eq.9:

$$y_{ci} = z_{ci} - x_{ci} \tag{9}$$

Suppose  $y_{ci}$  falls in the range  $[-sw_i \times a, sw_i \times b]$  following the probability density function  $g_{ci}(y_{ci})$ . Hence,  $x_{ci}$  locates in the range  $[sw_i \times (1-a), sw_i \times (1+b)]$ .

For two computers  $c_i$  and  $c_j$ , suppose that  $w_i < w_j$ . Then, the predicted execution time of a job with size s on  $c_i$  and  $c_j$  satisfy  $sw_i < sw_j$ . However, the range of the job's actual execution time on computer  $c_i$  is  $[sw_i \times (1-a), sw_i \times (1+b)]$  and may overlap with that on computer  $c_j$ , which is  $[sw_j \times (1-a), sw_j \times (1+b)]$ . Consequently, the actual execution time on computer  $c_i$  may be greater than that on  $c_j$ . In this case, the inaccurate predictions cause a misperception in the order of the actual execution times on these two computers. Depending on the individual scheduling algorithm, this misperception may lead to the wrong selection of a computer to which the job is sent. Similarly, the degree of misperception for a job with size s on two computers  $c_i$  and  $c_j$ , denoted by  $MD_c(c_i, c_j)$ , is defined by the probability that  $x_c \ge x_{cj}$  while  $z_{ci} < z_{cj}$ . This probability is denoted by  $P_r(x_{ci} \ge x_{cj} | z_{ci} < z_{cj})$ , which can be further transformed by Eq.10 with Eq.8 and 9.

$$P_{r}(x_{ci} \ge x_{cj} | z_{ci} < z_{cj}) = P_{r}(z_{ci} - y_{ci} \ge z_{cj} - y_{cj} | z_{ci} < z_{cj}) = P_{r}(y_{cj} \ge y_{ci} + sw_{j} - sw_{i} | w_{j} - w_{i} > 0)$$

$$(10)$$

That is,  $MD_c(c_i, c_j)$  is computed using the following equation:

$$MD_c(c_i, c_i) = P_r(v_{ci} \ge v_{ci} + sw_i - sw_i | w_i - w_i \ge 0)$$
 (11)

Applying a similar method to that used to compute  $MD(x_1, x_2)$ , the equation for computing  $MD_c(c_i, c_i)$  is expressed formally as follows.

$$MD_{c}(c_{i}, c_{j}) = \begin{cases} 0 & bsw_{j} \leq -asw_{i} + s(w_{j} - w_{i}) \\ \int_{-asw_{i}}^{bsw_{j} - s(w_{j} - w_{i})} \int_{y_{ci} + s(w_{j} - w_{i})}^{bsw_{j}} g_{ci}(y_{ci})g_{cj}(y_{cj})dy_{cj}dy_{ci} \\ -asw_{i} + s(w_{j} - w_{i}) < bsw_{j} < bsw_{i} + s(w_{j} - w_{i}) \end{cases}$$

$$\int_{-asw_{i}}^{bsw_{j}} \int_{y_{ci} + s(w_{j} - w_{i})}^{bsw_{j}} g_{ci}(y_{ci})g_{cj}(y_{cj})dy_{cj}dy_{ci} \quad bsw_{j} \geq bsw_{i} + s(w_{j} - w_{i})$$

$$(12)$$

The degree of misperception for a job with size s for n heterogeneous computers  $c_1, c_2, ..., c_n$ , denoted by  $\overline{MD_c}$ , is defined using the average of the degree of misperception for the job on any two computers, which can be computed using Eq.13.

$$\overline{MDc} = \frac{1}{C_n^2} \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} MDc(c_i, c_j)$$
(13)

# 3. An Evaluation of the Degree of Misperception

In the previous section the general formula for the degree of misperception for both job selection and resource allocation were presented. In this section an investigation is conducted as to how the parameters in these formulae impact on the value assigned to the degree of misperception.

#### 3.1 Job Selection

The parameters a and b represent the range of predicted errors in Eq.4-6. Fig.4.a and Fig.4.b show the impact of the parameters a and b on the value  $\overline{MD}$ . It is difficult to evaluate the impact of these parameters if the probability density function of predicted errors takes a general form. Therefore in the following parameter evaluation, the predicted error for the execution time of x is assumed to follow a uniform distribution in [-ax, bx], whose probability density function  $g_x(y_x)$  is expressed as:

$$g_x(y_x) = \frac{1}{(b+a)x}$$

Three types of job stream are investigated. The actual execution times in these job streams follow a uniform, Bounded Pareto and Exponential distribution, respectively. Their probability density functions f(x) are shown in Table 1. In the job stream following the exponential distribution, the job execution time has no upper limit. Only execution time values in [10, 14.6] are considered, as 99% of the execution times locate in this range according to its probability density function. This simplification does not impact on the accuracy of the results.

Table 1. The range of job execution times in the three job streams

	Uniform	Bounded Pareto	Exponential
f(x)	$\frac{1}{xu - xl}$	$\frac{\alpha \times x l^{\alpha}}{1 - (xl/xu)^{\alpha}} x^{-\alpha - 1} \ (\alpha = 1)$	$\frac{1}{\beta}e^{-(x-xl)/\beta} \ (\beta=1)$
[xl, xu]	[10, 100]	[10, 40]	[10, 14.6]

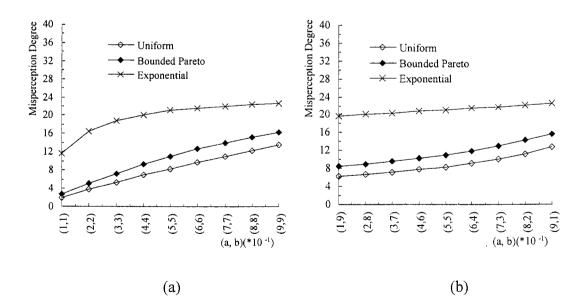


Figure.4. Impact of the parameters a and b on MD (a) the impact of the range size of predicted errors (b) the impact of the range location of predicted errors

In Fig.4.a, a and b increase from 10% to 90% with increments of 10%. This means that the range of predicted error for the actual execution time of x increases from [-0.1x, 0.1x] to [-0.9x, 0.9x], while the average predicted error remain unchanged (at 0).

As can be observed in Fig.4.a, under all three probability distributions, the degree of misperception increases as a and b increase. This is because as a and b increase the predicted execution times of jobs have a higher probability of overlapping with each other, which leads to the overall increase in  $\overline{MD}$ . This result suggests that when the average predicted error is the same, the range size of predicted errors is critical for the value of  $\overline{MD}$ .

It can also be observed that under the same a and b, the degree of misperception is highest under the exponential distribution, second highest under the Bound Pareto distribution and the lowest under the uniform distribution. This is because the size of the range of actual execution times is smallest when the execution times follow an exponential distribution and the largest when following a uniform distribution. This result suggests that the actual execution times will also influence the degree of misperception. This is demonstrated in Fig.5.

In Fig.4.b, the range size of the predicted error for the actual execution time of x remains unchanged (at x), while the location of the range shifts towards the left from [-0.1x, 0.9x] to [-0.9x, 0.1x]. The result of this is that the degree of misperception increases as the range location shifts leftwards (see Fig.4.b). The reason for this is as follows. Consider Fig.2.a and Fig.3.a. When (a, b) is (0.1, 0.9), the range of predicted errors for  $x_1$  and  $x_2$  are  $[-0.1x_1, 0.9x_1]$  and  $[-0.1x_2, 0.9x_2]$ , respectively. The size of the range where the predicted execution times for  $x_1$  and  $x_2$  overlap is

$$0.9x_1+0.1x_2-(x_2-x_1)$$

When (a, b) is (0.9, 0.1) however, the range of predicted errors are  $[-0.9x_1, 0.1x_1]$  and  $[-0.9x_2, 0.1x_2]$ , respectively, and the size of the overlapping ranges is

$$0.9x_2+0.1x_1-(x_2-x_1)$$

Since  $x_2$  is greater than  $x_1$ , hence,

$$0.9x_1+0.1x_2-(x_2-x_1)<0.9x_2+0.1x_1-(x_2-x_1)$$

this means that in general the size of the overlapping ranges is greater when (a, b) is (0.9, 0.1) than when (a, b) is (0.1, 0.9). This therefore leads to the increased degree of misperception. This result suggests that compared with an overestimation of execution time, the same level of underestimation may result in a higher degree of misperception.

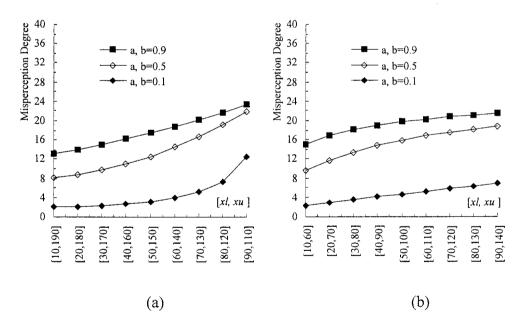


Figure.5. The impact of actual execution times on the degree of misperception (a) the impact of the range size of actual execution times (b) the impact of the range location of actual execution times

Fig. 5.a and Fig. 5.b show the impact of actual execution times on the degree of misperception. The results show the data for the actual execution times following a uniform distribution; the results for the Bounded Pareto distribution display a similar pattern. This study does not consider the execution times with an exponential distribution, since their range size (xu-xl) is fixed when 99% of the execution times are considered.

Fig.5.a shows the impact of the size of the range of actual execution times (i.e., xu-xl). In this same figure, the average of the actual execution times remains the same (at 100) while the range size of execution times decreases from 180 to 20 with decrements of 20. This experiment is conducted with different values of a and b. It can be observed in Fig.5.a that for the same values of a and b, the degree of misperception increases as the range size decreases. The reason for this is as follows: as the range size of the actual execution times decrease, the value of  $x_2-x_1$  (in Fig.2.a or Fig.3.a) decreases on average under the same a and b. As a result of this, the overlapping area of the two predicted execution times increases, which finally leads to an increase in  $\overline{MD}$ . This result suggests that when the average execution times are the same, a greater variance in execution time is of benefit, as this will reduce the degree of misperception.

Fig. 5.b demonstrates the impact of the location of the range of actual execution times. In Fig. 5.b, the range size of execution times remains constant (at 50) while the range shifts from [10, 60] to [90, 140]. As can be seen in Fig. 5.b, the degree of misperception increases in all cases as the range location shifts from [10, 60] to [90, 140]. The reason for this is that as the range location shifts, the mean execution time increases. Under the same a and b, the larger the actual execution time, the greater the range of its predicted error. Consequently, corresponding predicted execution times

have higher probability of overlapping with each other, which then incurs a higher degree of misperception. This result shows that when other parameters remain constant, the job stream with the greater average execution time tends to cause the highest degree of misperception.

#### 3.2 Resource allocation

In Eq.12 and 13, the parameters that influence  $\overline{MDc}$  include the error range parameters a and b, the computer weight  $w_i$  and the probability density function of predicted errors  $g_{ci}(y_{ci})$ . In the following experiments, the values of these parameters are given in Table 2 unless otherwise stated.

Table 2. The default values of the parameters for the experimentation

$w_I$	$w_{i}\text{-}w_{i-1} (2\leq i\leq n)$	n	S
10	5	6	50

In the following figures, the predicted error for the execution time of x is also assumed to follow a uniform distribution in  $[-asw_i, bsw_i]$ , whose probability density function  $g_{ci}(y_{ci})$  is expressed as follows:

$$g_{ci}(y_{ci}) = \frac{1}{(b+a)sw_i}$$

The parameters a and b indicate the range of the predicted error. Fig.6.a shows the impact of the range size on  $\overline{MDc}$ . The result has a similar pattern to that seen in Fig.4.a, which suggests that the range size of predicted errors is also critical for the value of  $\overline{MDc}$ . Similarly to Fig.4.b, Fig.6.b demonstrates the impact of the range location on  $\overline{MDc}$ . In the study of resource allocation,  $[sw_i(1-a), sw_i(1+b)]$  represents the range of actual execution times. Hence, the process where (a, b) shifts from (1, 9)

to (9, 1) means that the predicted execution time  $sw_i$  changes gradually from an underestimate to an overestimate. In Fig.6.b it can be seen that  $\overline{MD}_c$  decreases as (a, b) shifts from (1, 9) to (9, 1). These results coincide with those seen in Fig.4.b, in that compared with an overestimate of execution time, the same level of underestimation may incur a higher degree of misperception.

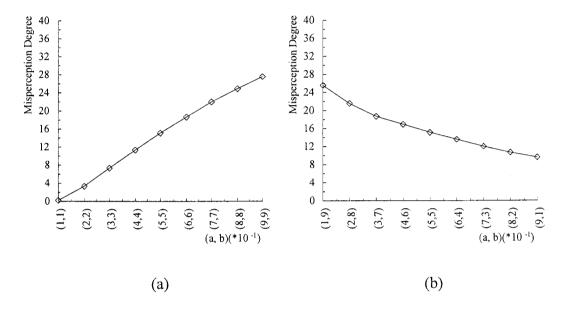


Figure.6. Impact of the parameters a and b on  $\overline{MD}_c$  (a) the impact of the range size of predicted errors (b) the impact of the range location of predicted errors

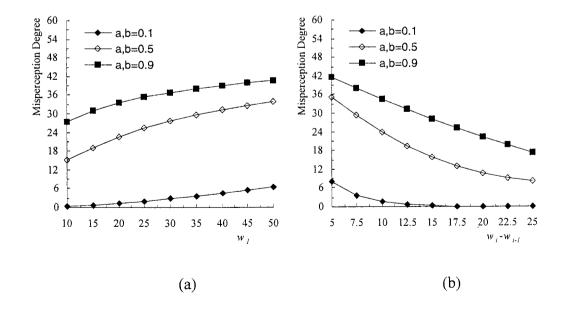


Figure.7. The impact of computer weight on  $\overline{MD_c}$  (a) the impact of the size of computer weights (b) the impact of the weight difference between computers

Fig.7.a shows the impact of computer weight on  $\overline{MD_c}$ . In Fig.7.a, the difference of the weight between computer  $c_i$  and  $c_{i-1}$  is fixed (at 5). As  $w_l$  increases (which means that resource  $c_l$  becomes slower), the weights of all the remaining computers increase accordingly. As can be observed in Fig.7.a,  $\overline{MD_c}$  increases as  $w_l$  increases under all values of a and b. This is because as  $w_l$  increases, the range of the actual execution time ( $[sw_l(1-a), sw_l(1+b)]$ ) also increases. This in turn increases the probability that the range of actual execution times on different computers overlap, which finally leads to an increased  $\overline{MD_c}$ . This result suggests that slower computers tend to generate higher degrees of misperception than faster computers.

Fig.7.b demonstrates the impact of weight difference between computers. In Fig.7.b, the difference between  $w_i$  and  $w_{i-1}$  ( $2 \le i \le n$ ) increases while the mean of these computer weights remains constant (at 70). It can be observed from this figure that  $\overline{MD}_c$  decreases as  $w_{i-1}w_{i-1}$  increases. This is because as  $w_{i-1}w_{i-1}$  increases, the difference between the predicted execution times on two computers  $c_i$  and  $c_j$  (i.e.,  $sw_j$ - $sw_i$ ) increases, which in turn reduces the probability that the ranges of their actual execution times overlap. This result suggests that resource pools with higher heterogeneity will achieve a lower degree of misperception.

#### 4. Conclusions

This paper investigates the underlying impact of inaccurate prediction on job selection and resource allocation. A new performance metric, termed the *degree of mis- perception*, is introduced in order to facilitate this exposition. General formulae have

been developed in order to calculate the degree of misperception for a variety of job streams and for distributed resource pools of varying levels of heterogeneity. The parameters that influence the degree of misperception are also investigated. This study will benefit the design and evaluation of different scheduling mechanisms for parallel and distributed systems that take prediction into account. It is likely that different scheduling policies will have different levels of sensitivity to this degree of misperception. Further work is planed to investigate how individual scheduling policies and specific performance measures are affected by this new performance metric.

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