How To Handle Missing Information Without Using NULL

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for Warwick University, CS253

SQL’s NULL Is A Disaster

See:
- Relational Database Writings 1985-1989 by C.J.Date with a special contribution by H.D. (as Andrew Warden)
- Relational Database Writings 1989-1991 by C.J.Date with Hugh Darwen
- Relational Database Writings 1991-1994 by C.J.Date
- Relational Database Writings 1994-1997 by C.J.Date
- Database Explorations by C.J. Date and Hugh Darwen (2010)

NULL

Cause of more debate and anguish than any other Fatal Flaw.
There’s even a split in the relational camp (E.F. Codd proposed “A-marks”, “I-marks” and a 4-valued logic).
There’s only one NULL. How many different reasons can there be for something being “missing”? Why NULL ruins everything –
- UNION of sets, cardinality of sets.
DeSTRUCTION of functional dependency theory SQL’s implementation of NULL is even worse than the best suggested by theoreticians. And it’s not completely BYPASSABLE, because SQL thinks that the sum of the empty set is NULL! Nor is it CORRECTABLE - the Shackel of Compatibility!

A Contradiction in Codd’s Proposed Treatment

“Every relation has at least one candidate key”
“Nulls aren’t permitted in the primary key”
“Nulls are permitted in alternate keys”
- Consider the relation resulting from the projection of PATEIN over RELIGION, a “nullable column”. Wait a moment! What’s this?
- List the candidate keys of this relation.
- Nominate the primary key.

Surprises Caused by SQL’s NULL

1. SELECT * FROM T WHERE X = Y
   UNION
   SELECT * FROM T WHERE NOT ( X = Y )
is not equal to SELECT * FROM T
2. SELECT SUM(X) + SUM(Y) FROM T
   is not equal to
   SELECT SUM(X + Y) FROM T
3. IF X = Y THEN ’Yes’; ELSE ’No’
is not equal to
   IF NOT ( X = Y ) THEN ’No’; ELSE ’Yes’

Why NULL Hurts Even More Than It Once Did

Suppose “x = x” returns Unknown
Can we safely conclude “x IS NULL”? Suppose x “is not the null value”?
Can we conclude “x IS NOT NULL”? Not in modern SQL!
How \( x = x \) Unknown Yet \( x \) NOT NULL

For example:

1. \( x \) is \texttt{ROW (1, null)} - or even \texttt{ROW(null, null)}
   \texttt{ROW(\ldots)} is a row “constructor”.

2. \( x \) is \texttt{POINT (1, null)}
   \texttt{POINT(a,b)} is a “constructor” for values in the user-defined data type \texttt{POINT}.

3. \( x \) is \texttt{ROW (\texttt{POINT(1,1)}, \texttt{POINT(null,3)})}

Consequences?

\textbf{x IS NULL (Case 1)}

What does \( x \) IS NULL MEAN? Think you know? Well, think again!

CREATE TABLE T ( C1 INT, C2 ROW ( F1 INT, F2 INT ) ) ;
INSERT INTO T VALUES ( NULL, NULL ) ;

<table>
<thead>
<tr>
<th>Query</th>
<th>Result Cardinality</th>
</tr>
</thead>
<tbody>
<tr>
<td>SELECT * FROM T WHERE C1 IS NULL</td>
<td>1</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE C2 IS NULL</td>
<td>1</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE ( C1, C1 ) IS NULL</td>
<td>1</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE ( C1, C2 ) IS NULL</td>
<td>1</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE ( C2, C2 ) IS NULL</td>
<td>1</td>
</tr>
</tbody>
</table>

But even this depends on our charitable interpretation of the ISO SQL standard.

\textbf{x IS NULL (Case 2)}

CREATE TABLE T ( C1 INT, C2 ROW ( F1 INT, F2 INT ) ) ;
INSERT INTO T VALUES ( NULL, ROW ( NULL, NULL ) ) ; -- note the difference from Case 1

<table>
<thead>
<tr>
<th>Query</th>
<th>Result Cardinality</th>
</tr>
</thead>
<tbody>
<tr>
<td>SELECT * FROM T WHERE C1 IS NULL</td>
<td>1</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE C2 IS NULL</td>
<td>1</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE ( C1, C2 ) IS NULL</td>
<td>0 !!!</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE ( C1, C1 ) IS NULL</td>
<td>1</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE ( C2, C2 ) IS NULL</td>
<td>0 !!!</td>
</tr>
</tbody>
</table>

\textbf{x IS NOT NULL}

So, what does \( x \) IS NOT NULL MEAN?

CREATE TABLE T ( C1 INT, C2 ROW ( F1 INT, F2 INT ) ) ;
INSERT INTO T VALUES ( NULL, ROW ( NULL, NULL ) ) ;

<table>
<thead>
<tr>
<th>Query</th>
<th>Result Cardinality</th>
</tr>
</thead>
<tbody>
<tr>
<td>SELECT * FROM T WHERE C1 IS NOT NULL</td>
<td>0</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE C2 IS NOT NULL</td>
<td>0</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE ( C1, C1 ) IS NOT NULL</td>
<td>0 !!!</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE ( C2, C2 ) IS NOT NULL</td>
<td>1 !!!</td>
</tr>
</tbody>
</table>

\textbf{Effects of Bad Language Design}

There are general language design lessons to be learned from this tangled web, as well as lessons about NULL:

- Enclosing an expression in parens should not change its meaning.
- Great caution is needed when considering pragmatic shorthands.
- All data types supported by a language should be “first-class”, for orthogonality.

\textbf{It Could Have Been Worse ...}

... if SQL had paid proper attention to degenerate cases.

SQL fails to recognise the existence of relations of degree zero (tables with no columns). These in turn depend on the existence of the 0-tuple. Suppose SQL had not made this oversight.

CREATE TABLE T ( C1 INT ) ;
INSERT INTO T VALUES ( ) ;

<table>
<thead>
<tr>
<th>Query</th>
<th>Result Cardinality</th>
</tr>
</thead>
<tbody>
<tr>
<td>SELECT * FROM T WHERE NOT ( C1 IS NULL )</td>
<td>1</td>
</tr>
<tr>
<td>SELECT * FROM T WHERE C1 IS NULL</td>
<td>1</td>
</tr>
</tbody>
</table>

\( C1 \) “is not the null value”; also, no field of \( C1 \) “is the null value”!

But it is also true that every field of \( C1 \) “is the null value”!
3-Valued Logic: The Real Culprit

Relational theory is founded on classical, 2-valued logic.

A relation \( r \) is interpreted as a representation of the extension of some predicate \( P \).

Let \( t \) be a tuple with the same heading as \( r \).

If tuple \( t \) is a member of \( r \), then the proposition \( P(t) \) is taken to be TRUE; otherwise (if \( t \) is not a member of \( r \)), \( P(t) \) is taken to be FALSE.

There is no middle ground. The Law of The Excluded Middle applies.

There is no way of representing that the truth of \( P(t) \) is unknown, or inapplicable, or otherwise concealed from us.

SQL’s WHERE clause arbitrarily splits at the TRUE/UNKNOWN divide.

Case Study Example

<table>
<thead>
<tr>
<th>Id</th>
<th>Name</th>
<th>Job</th>
<th>Salary</th>
</tr>
</thead>
<tbody>
<tr>
<td>1234</td>
<td>Anne</td>
<td>Lawyer</td>
<td>100,000</td>
</tr>
<tr>
<td>1235</td>
<td>Boris</td>
<td>Banker</td>
<td>?</td>
</tr>
<tr>
<td>1236</td>
<td>Cindy</td>
<td>?</td>
<td>70,000</td>
</tr>
<tr>
<td>1237</td>
<td>Davinder</td>
<td>?</td>
<td>?</td>
</tr>
</tbody>
</table>

Meaning (a predicate):
The person identified by Id is called Name and has the job of a Job, earning Salary pounds per year.

BUT WHAT DO THOSE QUESTION MARKS MEAN???

Summary of Proposed Solution

1. Database design:
   a. "vertical" decomposition
   b. "horizontal" decomposition

2. New constraint shorthands:
   a. "distributed key"
   b. "foreign distributed key"

3. New database updating construct:
   "multiple assignment"

4. Recomposition by query
to derive (an improved) PERS_INFO when needed

Database Design

a. "vertical" decomposition

Decompose into 2 or more relvars by projection
Also known as normalization.
Several degrees of normalization were described in the 1970s:
1NF, 2NF, 3NF, BCNF, 4NF, 5NF.
The ultimate degree, however, is 6NF: "irreducible relations".
(See "Temporal Data and The Relational Model", Date/Darwen/Lorentzos, 2003.)
A 6NF relvar consists of a key plus at most one other attribute.

Vertical Decomposition of PERS_INFO

<table>
<thead>
<tr>
<th>Id</th>
<th>Name</th>
<th>Id</th>
<th>Job</th>
<th>Id</th>
<th>Salary</th>
</tr>
</thead>
<tbody>
<tr>
<td>1234</td>
<td>Anne</td>
<td>1234</td>
<td>Lawyer</td>
<td>1234</td>
<td>100,000</td>
</tr>
<tr>
<td>1235</td>
<td>Boris</td>
<td>1235</td>
<td>Banker</td>
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<td>1236</td>
<td>?</td>
<td>1236</td>
<td>70,000</td>
</tr>
<tr>
<td>1237</td>
<td>Davinder</td>
<td>1237</td>
<td>?</td>
<td>1237</td>
<td>?</td>
</tr>
</tbody>
</table>

Meaning:
The person identified by Id is called Name.
The person identified by Id does the job of a Job.
The person identified by Id earns Salary pounds per year.

BUT WHAT DO THOSE QUESTION MARKS MEAN? (reprise)
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11 October, 2010

CS253: Topics in Databases 4

What We Have Achieved So Far

What started as a single table (PERS_INFO) is now a database (sub)schema (let’s call it PERS_INFO again), consisting of:

- CALLED (Id, Name)
- DOES_JOB (Id, Job)
- JOB_UNK (Id)
- UNEMPLOYED (Id)
- EARNS (Id, Salary)
- SALARY_UNK (Id)
- UNSALARIED (Id)

Next, we must consider the constraints needed to hold this design together (so to speak):

1. No two CALLED rows have the same Id. (Primary key)
2. Every row in CALLED has a matching row in either DOES_JOB, JOB_UNK, or UNEMPLOYED.
3. No row in DOES_JOB has a matching row in JOB_UNK.
4. No row in DOES_JOB has a matching row in UNEMPLOYED.
5. No row in JOB_UNK has a matching row in UNEMPLOYED.
6. Every row in DOES_JOB has a matching row in CALLED. (Foreign key)
7. Every row in JOB_UNK has a matching row in CALLED. (Foreign key)
8. Every row in UNEMPLOYED has a matching row in CALLED. (Foreign key)
9. Constraints 2 through 8 repeated, mutatis mutandis, for CALLED with respect to EARNS, SALARY_UNK and UNSALARIED.

WE NEED SOME NEW SHORTHANDS TO EXPRESS 2, 3, 4 AND 5.

Proposed Shorthands for Constraints

1. Id is a distributed key for (DOES_JOB, JOB_UNK, UNEMPLOYED). This addresses Constraints 3, 4 and 5.
2. Id is a distributed key for (EARNS, SALARY_UNK, UNSALARIED).
3. Id is a foreign distributed key in CALLED, referencing (DOES_JOB, JOB_UNK, UNEMPLOYED). This addresses Constraint 2.
4. Id is a foreign distributed key in CALLED, referencing (EARNS, SALARY_UNK, UNSALARIED).

Plus regular foreign keys in each of DOES_JOB, JOB_UNK, UNEMPLOYED, EARNS, SALARY_UNK, UNSALARIED, each referencing CALLED.

(Might also want UNEMPLOYED to imply UNSALARIED – how would that be expressed?)

So, now we have a schema and constraints. Next, how to add the data and subsequently update it? Are the regular INSERT/UPDATE/DELETE operators good enough?

Updating the Database: A Problem

How can we add the first row to any of our 7 tables?

Can’t add a row to CALLED unless there is a matching row in DOES_JOB, JOB_UNK or UNEMPLOYED and also a matching row in EARNS, SALARY_UNK or UNSALARIED.

Can’t add a row to DOES_JOB unless there is a matching row in CALLED.

Ditto JOB_UNK, UNEMPLOYED, EARNS, SALARY_UNK and UNSALARIED.

Impasse!
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### Updating the Database: Solution

"Multiple Assignment": doing several updating operations in a single "mouthful".

For example:

```sql
INSERT_TUPLE INTO CALLED { Id 1236, Name 'Cindy' } ,
INSERT_TUPLE INTO JOB_UNK { Id 1236 } ,
INSERT_TUPLE INTO EARNS { Id 1236, Salary 70000 } ;
```

Note very carefully the punctuation!

This triple operation is "atomic". Either it all works or none of it works.

Loosely speaking: operations are performed in the order given (to cater for the same target more than once), but intermediate states might be inconsistent and are not visible.

So, we now have a working database design. Now, what if the user wants to derive that original PERS_INFO table from this database?

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### To Derive PERS_INFO Relation from PERS_INFO Database

```sql
WITH (EXTEND JOB_UNK ADD ('Job unknown' AS Job_info)) AS T1,
(EXTEND UNEMPLOYED ADD ('Unemployed' AS Job_info)) AS T2,
(DOES_JOB RENAME (Job AS Job_info)) AS T3,
(EXTEND SALARY_UNK ADD ('Salary unknown' AS Sal_info)) AS T4,
(EXTEND UNSALARIED ADD ('Unsalaried' AS Sal_info)) AS T5,
(EXTEND EARNS ADD (CHAR(Salary) AS Sal_info)) AS T6,
(T6 { ALL BUT Salary }) AS T7,
(UNION ( T1, T2, T3 )) AS T8,
(UNION ( T4, T5, T7 )) AS T9,
JOIN ( CALLED, T8, T9 ) AS PERS_INFO :
```

Q.E.D.

---

### The New PERS_INFO

<table>
<thead>
<tr>
<th>Id</th>
<th>Name</th>
<th>Job_info</th>
<th>Sal_info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1234</td>
<td>Anne</td>
<td>Lawyer</td>
<td>100,000</td>
</tr>
<tr>
<td>1235</td>
<td>Boris</td>
<td>Banker</td>
<td>Salary unknown</td>
</tr>
<tr>
<td>1236</td>
<td>Cindy</td>
<td>Job unknown</td>
<td>70,000</td>
</tr>
<tr>
<td>1237</td>
<td>Davinder</td>
<td>Unemployed</td>
<td>Unsalaried</td>
</tr>
</tbody>
</table>

LOOK – NO QUESTION MARKS, NO NULLS!

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### How Much of All That Can Be Done Today?

- Vertical decomposition: can be done in SQL
- Horizontal decomposition: can be done in SQL
- Primary and foreign keys: can be done in SQL
- Distributed keys: can be done in (awful) longhand, if at all
- Foreign distributed keys can be done in (awful) longhand, if at all
- Multiple assignment: hasn’t caught the attention of SQL DBMS vendors, but Alphora’s D4 supports it.
- Recomposition query: can be done but likely to perform horribly. Might be preferable to store PERS_INFO as a single table under the covers, so that the tables resulting from decomposition can be implemented as mappings to that. But current technology doesn’t give clean separation of physical storage from logical design.

Perhaps something for the next generation of software engineers to grapple with?

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### Appendix A: Walk-through of Recomposition Query

We look at each step in turn, showing its effect.
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T1: EXTEND JOB_UNK ADD ('Job unknown' AS Job_info)

<table>
<thead>
<tr>
<th>Id</th>
<th>Job_info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1236</td>
<td>Job unknown</td>
</tr>
</tbody>
</table>

T2: EXTEND UNEMPLOYED ADD ('Unemployed' AS Job_info)

<table>
<thead>
<tr>
<th>Id</th>
<th>Job_info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1237</td>
<td>Unemployed</td>
</tr>
</tbody>
</table>

T3: DOES_JOB RENAME (Job AS Job_info)

<table>
<thead>
<tr>
<th>Id</th>
<th>Job_info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1234</td>
<td>Lawyer</td>
</tr>
<tr>
<td>1235</td>
<td>Banker</td>
</tr>
</tbody>
</table>

T4: EXTEND SALARY_UNK ADD ('Salary unknown' AS Sal_info)

<table>
<thead>
<tr>
<th>Id</th>
<th>Sal_info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1235</td>
<td>Salary unknown</td>
</tr>
</tbody>
</table>

T5: EXTEND UNSALARIED ADD ('Unsalaried' AS Sal_info)

<table>
<thead>
<tr>
<th>Id</th>
<th>Sal_info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1237</td>
<td>Unsalaried</td>
</tr>
</tbody>
</table>

T6: EXTEND EARNS ADD (CHAR(Salary) AS Sal_info)

<table>
<thead>
<tr>
<th>Id</th>
<th>Salary</th>
<th>Sal_info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1234</td>
<td>100,000</td>
<td></td>
</tr>
<tr>
<td>1235</td>
<td>70,000</td>
<td></td>
</tr>
</tbody>
</table>

Salary and Sal_info differ in type. Sal_info contains character representations of Salary values (hence left justified).
**HACD.1 How To Handle Missing Information Without Using NULL**

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<table>
<thead>
<tr>
<th>Id</th>
<th>Salary</th>
<th>Sal_info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1234</td>
<td>100,000</td>
<td>100,000</td>
</tr>
<tr>
<td>1236</td>
<td>70,000</td>
<td>70,000</td>
</tr>
</tbody>
</table>

**T7: T6 (ALL BUT Salary)**

<table>
<thead>
<tr>
<th>Id</th>
<th>Job_info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1234</td>
<td>Job unknown</td>
</tr>
<tr>
<td>1236</td>
<td>Lawyer</td>
</tr>
<tr>
<td>1235</td>
<td>Banker</td>
</tr>
<tr>
<td>1236</td>
<td>Job unknown</td>
</tr>
<tr>
<td>1237</td>
<td>Unemployed</td>
</tr>
</tbody>
</table>

**T9: UNION (T4, T5, T7)**

<table>
<thead>
<tr>
<th>Id</th>
<th>Name</th>
<th>Id</th>
<th>Job_info</th>
<th>Id</th>
<th>Sal_info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1234</td>
<td>Anne</td>
<td>1234</td>
<td>Lawyer</td>
<td>1234</td>
<td>100,000</td>
</tr>
<tr>
<td>1235</td>
<td>Boris</td>
<td>1235</td>
<td>Banker</td>
<td>1235</td>
<td>Salary unknown</td>
</tr>
<tr>
<td>1236</td>
<td>Cindy</td>
<td>1236</td>
<td>Job unknown</td>
<td>1236</td>
<td>70,000</td>
</tr>
<tr>
<td>1237</td>
<td>Davinder</td>
<td>1237</td>
<td>Unemployed</td>
<td>1237</td>
<td>Unsalaried</td>
</tr>
</tbody>
</table>

**PERS_INFO: JOIN (CALLED, T8, T9)**

<table>
<thead>
<tr>
<th>Id</th>
<th>Name</th>
<th>Id</th>
<th>Job_info</th>
<th>Id</th>
<th>Sal_info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1234</td>
<td>Anne</td>
<td>1234</td>
<td>Lawyer</td>
<td>1234</td>
<td>100,000</td>
</tr>
<tr>
<td>1235</td>
<td>Boris</td>
<td>1235</td>
<td>Banker</td>
<td>1235</td>
<td>Salary unknown</td>
</tr>
<tr>
<td>1236</td>
<td>Cindy</td>
<td>1236</td>
<td>Job unknown</td>
<td>1236</td>
<td>70,000</td>
</tr>
<tr>
<td>1237</td>
<td>Davinder</td>
<td>1237</td>
<td>Unemployed</td>
<td>1237</td>
<td>Unsalaried</td>
</tr>
</tbody>
</table>

The Very End