

Time allowed: 2 hours.

Answer **FOUR** questions.

Read carefully the instructions on the answer book and make sure that the particulars required are entered on each answer booklet.

1. (a) State the natural deduction rules: \rightarrow -introduction, \rightarrow -elimination, \vee -elimination, \neg -introduction, \neg -elimination. [8]
(b) Prove $p \rightarrow (q \vee r)$, $\neg q \vdash \neg r \rightarrow \neg p$ by natural deduction. [8]
(c) Describe a procedure $\text{NNF}(\phi)$ which, given an implication-free propositional formula ϕ , returns a negation normal form for ϕ . [9]

 2. (a) What does it mean for the set of rules of a deductive system to be *sound* and to be *complete*? [5]
(b) Sketch a proof that the natural deduction rules are complete for propositional logic. [9]
(c) Is natural deduction sound and complete for predicate logic? [3]
(d) What does it mean to say that the algorithmic problem of determining the validity of a predicate logic formula is *undecidable*? [4]
(e) What is the practical significance for automated reasoning of soundness, completeness and undecidability results in logic? [4]
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3. (a) Write down a formula ϕ which is equivalent to $S(x) \wedge \forall y(P(x) \rightarrow Q(y))$, and such that $f(y, y)$ is free for x in ϕ . [8]
- (b) Explain what is meant by each of the following:
- i. a *model* of predicate logic;
 - ii. an *environment* with respect to a model;
 - iii. that ϕ_1, \dots, ϕ_n *semantically entail* ψ ;
 - iv. that ϕ_1, \dots, ϕ_n are *consistent*.
- [8]
- (c) Define $f^{\mathcal{M}}$ so that, together with the following definitions:

$$A = \{a, b, c, d\}$$

$$P^{\mathcal{M}} = \{(a, b), (b, c), (c, d), (d, a)\}$$

we obtain a model \mathcal{M} in which the formula $\neg \exists x P(f(x), x)$ is satisfied. [9]

4. (a) What is meant by satisfaction of a Hoare triple $\langle \phi \rangle P \langle \psi \rangle$ under partial correctness? What is meant by satisfaction of it under total correctness? [8]
- (b) State and briefly explain each of the following proof rules for partial correctness: **Composition**, **Assignment**, **If-statement**, **Partial-while**, **Implied**. [8]
- (c) Let Div denote the following code:

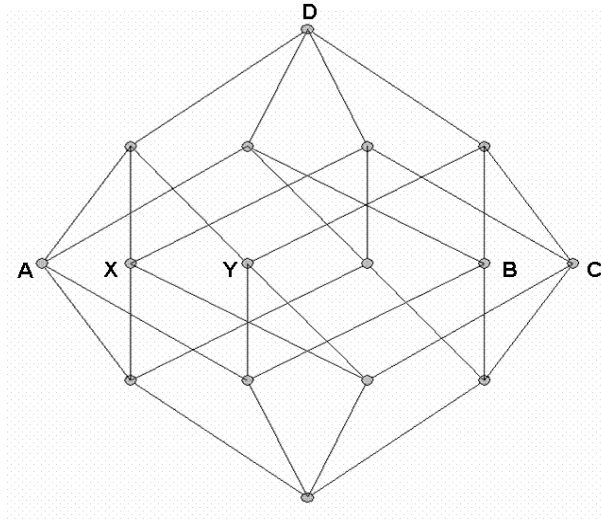
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z = 0;
while (x >= y) {
    z = z + 1;
    x = x - y;
}

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Show $\vdash_{tot} \langle x = x_0 \wedge x \geq 0 \rangle \text{Div} \langle x_0 = y \cdot z + x \wedge x < y \wedge x \geq 0 \rangle$ by a proof tableau. [9]

5. (a) A *Boolean algebra* is a distributive lattice in which every element has a complement. Express this characterisation formally by explicitly stating the axioms that define a Boolean algebra. [5]
- (b) Prove that every element p in a Boolean algebra has a unique complement [5]



- (c) Suppose that the above figure is interpreted as a Hasse diagram representing the set of subsets of $\{1, 2, 3, 4\}$ ordered by inclusion, and that the points X and Y represent the subsets $\{1, 3\}$ and $\{2, 3\}$ respectively. Identify the subsets that are represented by the points A, B, C and D respectively. [5]
- (d) Suppose that the Hasse diagram is interpreted as representing the set of logically distinct propositional formulas that can be constructed from atomic propositions x and y ordered by the relation $\phi \leq \psi$ if and only if $\phi \rightarrow \psi$, and that x and y are represented by the points X and Y respectively. Identify four propositions that are represented by the points A, B, C and D respectively. [5]
- (e) Explain informally how you would use the Alloy tool to search for Boolean algebras with at most 8 elements. [5]